

Overview of State-Dependent Importance Sampling for Simulating Jackson Networks

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Outline

- Motivation
- Model and Notation
- MAP-Based Approach
(Kroese and Nicola 1999)
- Adaptive IS Heuristics
(Rubinstein 1999, de Boer and Nicola 2002)
- State-Dependent Heuristics
(Nicola and Zaburnenko 2004)
- Game-Based Approach
(Dupuis, Sezer and Wang 2007)
- Conclusions and Further Research

Motivation

- Efficient simulation of rare events in queueing models (e.g., buffer overflow)
- Applications
 - Computer Systems: performance/dependability evaluation
 - Telecommunication Networks: QoS management and control
 - Manufacturing: scheduling and resource allocation
- Importance sampling (IS) technique
 - State-independent (static) IS (does not always work!)
 - state-dependent (dynamic) IS (generally more effective!)
- Brief overview of recent advances in dynamic IS

Model and Notation

- Jackson queueing network with n nodes

λ_i external arrival rate at Node i

μ_i service rate at Node i

p_{ij} routing probability from i to j

p_{ie} routing probability from i to *exit*

Λ_i total arrival rate at Node i

(from traffic equations):

$$\Lambda_i = \lambda_i + \sum_{\forall j} \Lambda_j p_{ji}, \quad i = 1, \dots, n$$

$\Lambda_i < \mu_i, \quad i = 1, \dots, n$ (all queues are stable!)

Model and Notation (contd.)

- Notation

- $X_{i,t}$ ($i = 1, \dots, n$) number at Node i at $t \geq 0$
- $\mathbf{X}_t = (X_{1,t}, \dots, X_{n,t})$, $t \geq 0$, Markov process
- $S_t = \sum_{\forall i} \beta_i X_{i,t}$, with $\beta_i = \{0, 1\}$, $i = 1, \dots, n$

When $\forall i \beta_i = 1$, $S_t \equiv$ network population

When only $\beta_i = 1$, $S_t \equiv$ buffer i content

- Starting from $S_0 = i \geq 1$

τ_K first time $S_t = K$

τ_0 first time $S_t = 0$

- rare event of interest:

S_t hits K before return to 0

- associated probability: $\gamma_i(K) = \mathbb{P}(\tau_K < \tau_0)$

State-Independent (Static) IS

- heuristics based on large deviations, e.g.,
(Parekh and Walrand 1989)
(Frater et al. 1989, 1991 and 1994)
(Juneja and Nicola 2005)
- heuristics based on effective bandwidth, e.g.,
(Chang et al. 1994)
(De Veciana et al. 1994)
(L'Ecuyer and Champoux 2001)
- not always asymptotically efficient!
(Glasserman et al. 1995 and 1997)
(Randhawa and Juneja 2004)
(de Boer 2004)

State-Dependent (Dynamic) IS

- Two basic ideas:
 - approximate the zero-variance C.O.M.
 - “push” the system close to the most likely path to the rare set (MLP C.O.M.)
- Zero-variance and MLP C.O.M.s:
 - not known (require the true probability!)
 - state-dependent (even for the $M/M/1$ queue!)
 - strong dependence close to boundaries (when one or more queues are empty!)
 - weak or no dependence in the interior
- Challenges:
 - determine C.O.M.s (at boundaries/interior)
 - appropriately combine these C.O.M.s
 - dependence range along boundaries is crucial for asymptotic efficiency!
 - boundaries increase with number of queues

State-Dependent (Dynamic) IS (contd.)

- Generally more effective than static IS!
- Formal MAP approach
(Kroese and Nicola 2002)
 - for 2-node tandem networks
 - provably effective
 - difficult for networks with many nodes?
- Adaptive IS heuristics
 - cross-entropy (CE) (pioneered by Rubinstein)
 - CE for QN models (de Boer and Nicola 2002)
 - stoch. approximation (Ahamed et al. 2005)
 - less effective for networks with many nodes:
excessive computational effort;
convergence problems?

State-Dependent (Dynamic) IS (contd.)

- State-dependent IS heuristics
(Nicola and Zaburnenko 2007)
 - less computational effort than adaptive IS
 - effective for tandem/parallel topologies
 - not developed for general Jackson networks?
- Formal game-theoretic approach
(Dupuis and Wang 2008)
 - rigorous optimal control framework
 - provably effective
 - implementation and performance problems for networks with many nodes?

MAP-Based Approach

- First (formal) approach to state-dependent IS (Kroese and Nicola 1999, 2002)
- Stable two-node tandem network (λ, μ_1, μ_2)
- Rare event: starting from $(X_{1,0} = i, X_{2,0} = 1)$, second buffer $(X_{2,t})$ hits some high level K before returning to 0; associated prob. $\gamma_i(K)$
- $(X_{1,t})$: birth-death process (λ, μ_1)
 (D_t) : departure process from Node 1
 (E_t) : pure death process (rate μ_2)
- Define $S_t \stackrel{\text{def}}{=} X_{2,0} + (D_t - E_t)$, $t \geq 0$,
then $S_t \equiv X_{2,t}$, $0 \leq t \leq \min\{\tau_0, \tau_K\}$

MAP-Based Approach (contd.)

- $(X_{1,t}, S_t)$: Markov additive process (MAP)

$(X_{1,t})$: modulating (driving) process;
continuous-time birth-death process (λ, μ_1)

(S_t) : modulated (additive) process,
with increments D_t and death rate μ_2

- $M_t(\theta)$: matrix of moment gen. functions

$$(M_t(\theta))_{i,j} \stackrel{\text{def}}{=} \mathbf{E}_i e^{\theta S_t} I_{\{X_t=j\}}, \quad 0 \leq i, j \leq \infty$$

- Then $M_t(\theta) = e^{tG(\theta)}$, $t \geq 0$, with $G(\theta)$ given by

$$\begin{pmatrix} -\lambda - \mu_2 + \mu_2 e^{-\theta} & \lambda & & & \\ \mu_1 e^{\theta} & -\lambda - \mu_1 - \mu_2 + \mu_2 e^{-\theta} & \lambda & & \\ & \ddots & \ddots & \ddots & \\ & & & \mu_1 e^{\theta} & -\mu_1 - \mu_2 + \mu_2 e^{-\theta} \end{pmatrix}$$

MAP-Based Approach (contd.)

- $\kappa(\theta)$: largest positive eigen value of $G(\theta)$
- $\mathbf{w}(\theta) = \{w_i(\theta), i = 0, 1, \dots\}$:
right-eigenvector corresponding to $\kappa(\theta)$
- Exponential C.O.M.: MAP $(\tilde{X}_{1,t}, \tilde{S}_t)$ given by
 $(\tilde{X}_{1,t})$: modulating process; a continuous-time
birth-death process with conjugate rates

$$\tilde{\lambda}(i) = G_{i,i+1}(\theta) \frac{w_{i+1}(\theta)}{w_i(\theta)}, \quad i = 0, 1, \dots$$

$$\tilde{\mu}_1(i) = G_{i,i-1}(\theta) \frac{w_{i-1}(\theta)}{w_i(\theta)}, \quad i = 1, 2, \dots$$

rates depend on i (content of the first buffer)

- (\tilde{S}_t) : modulated additive process, with
increments \tilde{D}_t and death rate $\tilde{\mu}_2 = \mu_2 e^{-\theta}$

MAP-Based Approach (contd.)

- To estimate $\gamma_i(K)$, perform IS simulation with the optimal tilting parameter θ^* , as determined from $\kappa(\theta^*) = 0$
(Asmussen and Rubinstein 1995)
- IS estimator is provably asymptotically efficient!
- MAP-based approach application to larger networks seems difficult?
- Conclusion: more research on state-dep. IS!

Adaptive IS Heuristics

- The method of **Cross-Entropy (CE)** for rare event simulation (**Rubinstein 1999, de Boer and Nicola 2000, 2002**)
- The CE method is an **iterative procedure**; each iteration involves two basic steps:
 1. **generate** samples according to a given **prob. measure** (specified by a set of parameters)
 2. **update** parameters of the **prob. measure** (based on the samples collected in Step 1) to produce “better” samples in the next iteration
- The goal is to **converge** to a **prob. measure** sufficiently close to the **zero-variance C.O.M.**

Adaptive IS Heuristics (contd.)

- Consider the simulation of an $M/M/1$ queue
 - original prob. measure $f(t; \mathbf{v})$:
 $\mathbf{v} = (\lambda, \mu)$ (with $\lambda + \mu = 1$)
 - for overflow level K , we wish to estimate $\gamma(K)$ using importance sampling
 - change of measure $f(t; \hat{\mathbf{v}})$:
 $\hat{\mathbf{v}} = (\hat{\lambda}, \hat{\mu})$ (also, $\hat{\lambda} + \hat{\mu} = 1$)
- The CE method is used to determine an optimal parameter vector $\hat{\mathbf{v}}^*$

Adaptive IS Heuristics (contd.)

- The CE algorithm (to determine $\hat{\mathbf{v}}^*$):
 1. Set $\hat{\mathbf{v}}_0 = \mathbf{v}$ and $j = 1$ (iteration counter)
 2. Update $\hat{\mathbf{v}}_j$:
 - generate cycles w_1, \dots, w_m using $f(t; \hat{\mathbf{v}}_{j-1})$
 - select the “best” 1% cycles (with highest levels). Set $K_j \leq K$ to the lowest level reached
 - update $\hat{\mathbf{v}}_j = (\hat{\lambda}_j, \hat{\mu}_j)$ (CE minimization):
$$\hat{\lambda}_j = \frac{\sum_{i=1}^m I_i(K_j) L_i \times A_{ij}}{\sum_{i=1}^m I_i(K_j) L_i \times (A_{ij} + D_{ij})}, \quad \hat{\mu}_j = 1 - \hat{\lambda}_j$$

A_{ij} : arrivals in cycle i before reaching K_j

D_{ij} : departures in cycle i before reaching K_j
 3. if $K_j = K$, set $\hat{\mathbf{v}}^* = \hat{\mathbf{v}}_j$ and go to Step 4; otherwise, set $j = j + 1$ and go to Step 2
 4. use $\hat{\mathbf{v}}^*$ to estimate $\gamma(K)$ via IS

Adaptive IS Heuristics (contd.)

- The CE algorithm can be used to optimize a state-dependent C.O.M. $\underline{\hat{v}} = (\underline{\hat{\lambda}}, \underline{\hat{\mu}})$ with

$$\underline{\hat{\lambda}} = (\hat{\lambda}(0), \dots, \hat{\lambda}(b))$$

$$\underline{\hat{\mu}} = (\hat{\mu}(1), \dots, \hat{\mu}(b))$$

b is the dependence range along boundary
(generally critical for asymptotic efficiency!)

- Adaptive IS:
 - robust and effective for small networks
 - less effective for large networks:
excessive computational effort
convergence problems?
- Related work based on stoch. approximation
(Ahamed, Borkar and Juneja 2005)

State-Dependent Heuristics

- Goal: “push” the system close to the most likely path (MLP) to the rare set (Nicola and Zaburdenko 2007)
- Methodology:
 - approximate the MLP C.O.M. along a few important boundaries and in the interior of the state-space (e.g., via time-reversal arguments)
 - appropriately combine these C.O.M.s (e.g., via simple linear interpolation)
 - set/adjust the dependence range (boundary thickness) to attain asymptotic efficiency

State-Dependent Heuristics (contd.)

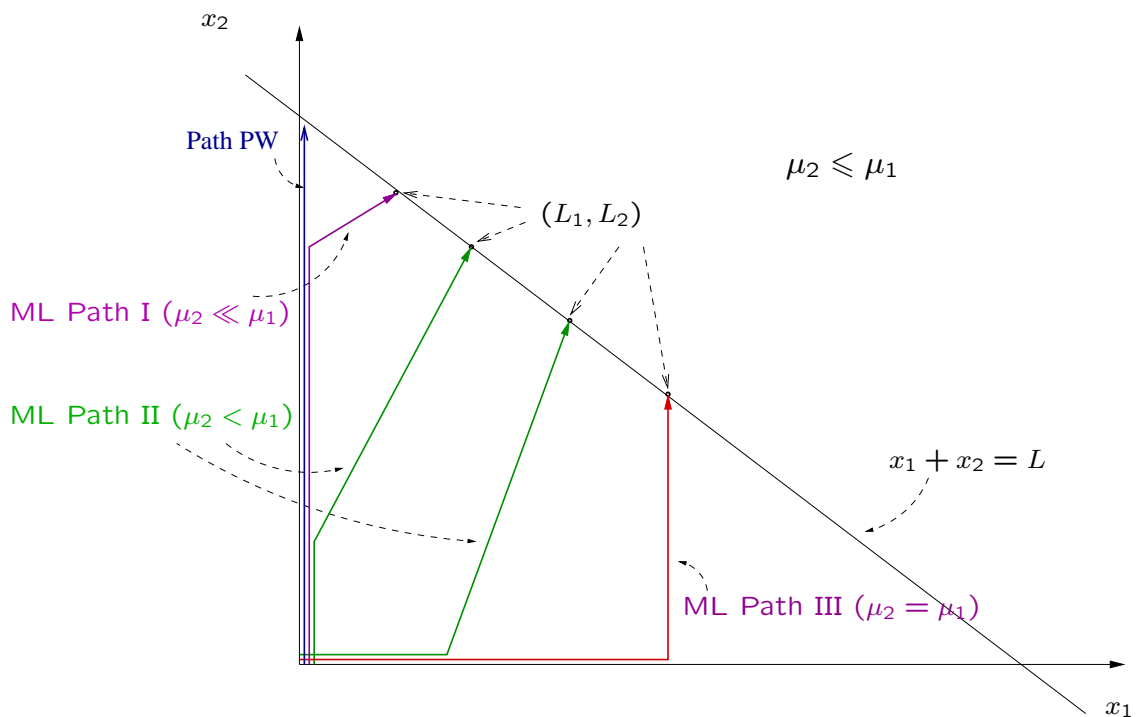
Time reversal of the 2-node tandem network



a) Forward time process

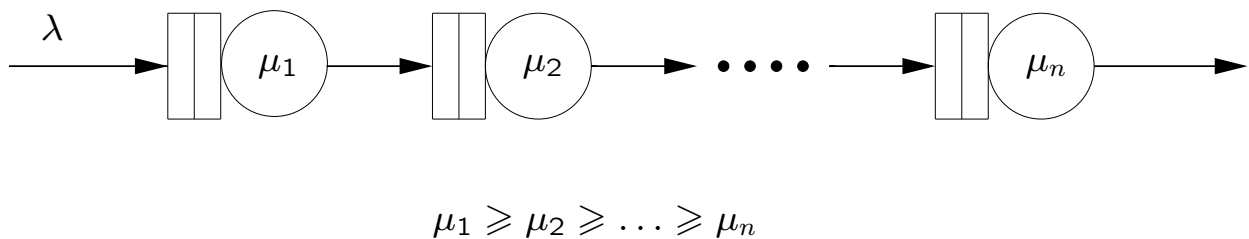


b) Reverse time process



State-Dependent Heuristics (contd.)

SDH for n -node tandem networks



From an empty network:

- initially ($x_i = 0, \forall i$), no change of measure;
- as $x_1 \geq 1$, start 'pushing' Node 1;
- as $x_2 \geq 1$, gradually go to 'pushing' Node 2;
- as $x_3 \geq 1$, gradually go to 'pushing' Node 3;
- etc.

State-Dependent Heuristics (contd.)

- Advantages:
 - relatively **simple** to develop and implement
 - **no computational effort** to determine C.O.M.
 - empirical results show **effectiveness** for
tandem/parallel networks
with any number of nodes
feed-forward/feedback networks
of small and moderate size
- Challenges:
 - approximating **MLP C.O.M. along boundaries**
 - guidelines to set the **dependence range**
 - proof of **asymptotic efficiency**

Game-based Approach

- A formal **game-theoretic** foundation for the development of provably efficient state-dependent IS schemes (Dupuis, Sezer and Wang 2007)
- Goal: “push” the system along the **most likely path (MLP)** to the rare set
- Methodology:
 - **stochastic optimal control** formulation to minimize the variance of the IS estimator (in the limit as $K \rightarrow \infty$, converges to a **deterministic optimal control** problem)
 - **approximate solution** to the associated **DPE** provides “key ingredients” to determine **MLP C.O.M.s** along distinct boundaries
 - a proper **weighted sum** of these C.O.M.s yields an **asymptotically efficient** IS scheme

Game-based Approach (contd.)

- Advantage:
 - rigorous and systematic framework for the construction of asymptotically efficient IS schemes for Jackson networks
 - simultaneous estimation of probabilities for different overflow events
- Challenges:
 - despite asymptotic efficiency, performance is sensitive to boundary thickness; rel. error may grow quickly (more than linear) with K
 - number of boundaries $(2^n - 1)$ increases exponentially with the number of nodes (n)
 - implementation and performance issues for large networks

Conclusions and Further Research

- Significant and promising recent advances
 - Adaptive IS heuristics
 - State-dependent IS heuristics
 - rigorous control-theoretic approach
- Further research
 - formal study of boundary thickness' impact on asymptotic efficiency and actual performance
 - effectiveness for large networks (applicability and performance issues)
 - reversibility-based approach (work in progress)
 - extensions to non-Jackson queueing networks